

Unit 1.4 Mathematical Backgrounds

Algorithms

EE/NTHU

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Mathematical Backgrounds

Monotonicity

- A function $f(n)$ is **monotonically increasing** if $m \leq n$ implies $f(m) \leq f(n)$.
- A function $f(n)$ is **monotonically decreasing** if $m \leq n$ implies $f(m) \geq f(n)$.
- A function $f(n)$ is **strictly increasing** if $m < n$ implies $f(m) < f(n)$.
- A function $f(n)$ is **strictly decreasing** if $m < n$ implies $f(m) > f(n)$.

Floor and ceiling functions

- For any real number x , we denote the greatest integer less than or equal to x by $\lfloor x \rfloor$ and the least integer greater than or equal to x by $\lceil x \rceil$.

- For any real x

$$x - 1 < \lfloor x \rfloor \leq x \leq \lceil x \rceil < x + 1. \quad (1.4.1)$$

- For any integer n ,

$$\lceil n/2 \rceil + \lfloor n/2 \rfloor = n. \quad (1.4.2)$$

- For any real number $x \geq 0$ and integers $m, n > 0$,

$$\lceil \lceil x/m \rceil / n \rceil = \lceil x/(mn) \rceil, \quad (1.4.3)$$

$$\lfloor \lfloor x/m \rfloor / n \rfloor = \lfloor x/(mn) \rfloor, \quad (1.4.4)$$

$$\lceil m/n \rceil \leq (m + (n - 1))/n, \quad (1.4.5)$$

$$\lfloor m/n \rfloor \leq (m + (n - 1))/n. \quad (1.4.6)$$

- The floor function $\lfloor x \rfloor$ is monotonically increasing, so is the ceiling function $\lceil x \rceil$.

Mathematical Backgrounds, II

Modular arithmetic

- For any integer m and positive integer n , the value $m \bmod n$ is the remainder (or residue) of the quotient m/n :

$$m \bmod n = m - \lfloor m/n \rfloor n. \quad (1.4.7)$$

- If $(a \bmod n) = (b \bmod n)$, we write $a \equiv b \pmod{n}$ and say a is equivalent to b , modulo n .
- $a \equiv b \pmod{n}$ if a and b have the same remainder when divided by n .
- $a \equiv b \pmod{n}$ if and only if n is a divisor of $b - a$.
- We write $a \not\equiv b \pmod{n}$ if a is not equivalent to b , modulo n .

Mathematical Backgrounds, III

Polynomials

- Given a nonnegative integer n , a polynomial in x of degree n is a function $p(x)$ of the form

$$p(x) = \sum_{k=0}^n a_k x^k, \quad (1.4.8)$$

where the constants a_0, a_1, \dots, a_n are the coefficients of the polynomial and $a_n \neq 0$.

- A polynomial is asymptotically positive if and only if $a_n > 0$.
- For an asymptotically positive polynomial $p(x)$ of degree n , we have $p(x) = \Theta(x^n)$.
- For any real constant $c \geq 0$ then function x^c is monotonically increasing, and for any real constant $c \leq 0$, the function x^c is monotonically decreasing.
- We say that a function $f(x)$ is polynomial bounded if $f(x) = \mathcal{O}(x^k)$ for some constant k .

Mathematical Backgrounds, IV

Exponentials

- For all real $a > 0$, m and n , we have the following identities:

$$a^0 = 1, \quad (1.4.9)$$

$$a^1 = a, \quad (1.4.10)$$

$$a^{-1} = 1/a, \quad (1.4.11)$$

$$(a^m)^n = a^{mn}, \quad (1.4.12)$$

$$(a^m)^n = (a^n)^m, \quad (1.4.13)$$

$$a^m \cdot a^n = a^{m+n}. \quad (1.4.14)$$

- For all n and $a \geq 1$, the function a^n is monotonically increasing in n . When convenient, we assume $0^0 = 1$.
- For all real constants a and b such that $a > 1$,

$$\lim_{n \rightarrow \infty} \frac{n^b}{a^n} = 0, \quad (1.4.15)$$

thus

$$n^b = o(a^n). \quad (1.4.16)$$

That is any exponential function with a base strictly greater than 1 grows faster than any polynomial function.

Mathematical Backgrounds, V

- Let e be the base of the natural logarithm function, $e = 2.71828 \dots$, we have for all real x

$$e^x = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots = \sum_{k=0}^{\infty} \frac{x^k}{k!}. \quad (1.4.17)$$

- For all real x , we have the inequality

$$e^x \geq 1 + x, \quad (1.4.18)$$

with the equality holds only when $x = 0$.

- When $|x| \leq 1$, we have the approximation

$$1 + x \leq e^x \leq 1 + x + x^2. \quad (1.4.19)$$

- Considering $x \rightarrow 0$, we have

$$e^x = 1 + x + \Theta(x^2). \quad (1.4.20)$$

- For all real x , we have

$$\lim_{n \rightarrow \infty} \left(1 + \frac{x}{n}\right)^n = e^x. \quad (1.4.21)$$

Mathematical Backgrounds, VI

Logarithms

- The following notations are adopted

$$\begin{aligned} \lg n &= \log_2 n && \text{(binary logarithm),} \\ \ln n &= \log_e n && \text{(natural logarithm),} \\ \lg^k n &= (\lg n)^k && \text{(exponentiation),} \\ \lg \lg n &= \lg(\lg n) && \text{(composition).} \end{aligned}$$

- We also adopt the convention that the logarithm functions only apply to the next term in the formula, so that $\lg n + k = (\lg n) + k$.
- If $b > 1$ is a constant, then for $n > 0$ the function $\log_b n$ is strictly increasing.
- For all $a > 0$, $b > 0$, $c > 0$ and n ,

$$\begin{aligned} a &= b^{\log_b a}, \\ \log_c(ab) &= \log_c a + \log_c b, \\ \log_b a^n &= n \log_b a, \\ \log_b a &= \frac{\log_c a}{\log_c b}, \\ \log_b(1/a) &= -\log_b a, \\ \log_b a &= \frac{1}{\log_a b}, \\ a^{\log_b c} &= c^{\log_b a}, \end{aligned}$$

where the base of each logarithm is not 1.

Mathematical Backgrounds, VII

- When $|x| < 1$,

$$\ln(1+x) = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \frac{x^5}{5} - \dots \quad (1.4.22)$$

- For $x > -1$,

$$\frac{x}{1+x} \leq \ln(1+x) \leq x. \quad (1.4.23)$$

where the equality holds only for $x = 0$.

- A function $f(n)$ is **polylogarithmically bounded** if $f(n) = \mathcal{O}(\lg^k n)$ for some constant k . Since

$$\lim_{n \rightarrow \infty} \frac{\lg^b n}{(2^a)^{\lg n}} = \lim_{n \rightarrow \infty} \frac{\lg^b n}{n^a} = 0, \quad (1.4.24)$$

we have

$$\lg^b n = o(n^a) \quad (1.4.25)$$

for any constant $a > 0$. Thus, any positive polynomial function grows faster than any polylogarithmic function.

- Change the base of a logarithm from one constant to another changes the value by a constant factor, so in conjunction with the \mathcal{O} -notation, the use of \log , or \lg or \log_2 are equivalent.

Mathematical Backgrounds, VIII

Factorials

- The factorial function, $n!$, is defined for integers $n \geq 0$ as

$$n! = \begin{cases} 1 & \text{if } n = 0, \\ n \cdot (n-1)! & \text{if } n > 0. \end{cases} \quad (1.4.26)$$

Thus, $n! = 1 \cdot 2 \cdot 3 \cdots n$.

- A weak upper bound on the factorial function is $n! \leq n^n$.
- Stirling's approximation**

$$n! = \sqrt{2\pi n} \left(\frac{n}{e}\right)^n \left(1 + \Theta\left(\frac{1}{n}\right)\right). \quad (1.4.27)$$

- Thus

$$\begin{aligned} n! &= o(n^n), \\ n! &= \omega(2^n), \\ \lg(n!) &= \Theta(n \lg n). \end{aligned}$$

- For $n \geq 1$

$$n! = \sqrt{2\pi n} \left(\frac{n}{e}\right)^n e^{\alpha_n}, \quad (1.4.28)$$

where

$$\frac{1}{12n+1} < \alpha_n < \frac{1}{12n}.$$

Mathematical Backgrounds, IX

Function iteration

- The notation $f^{(i)}(x)$ is used to denote function $f(x)$ iteratively applied i times to an initial value of x .
- That is, let $f(x)$ be a function over the reals. Given a nonnegative integer i , define

$$f^{(i)}(x) = \begin{cases} x & \text{if } i = 0, \\ f(f^{(i-1)}(x)) & \text{if } i > 0. \end{cases} \quad (1.4.29)$$

- For example, if $f(x) = 2x$, then $f^{(i)}(x) = 2^i x$.
- Note the difference of $f^{(i)}(x)$ and $f^i(x)$, which is $f(x)$ raised to the i th power.

Iterative logarithm function

- The **iterative logarithm function** is defined as

$$\lg^* x = \min\{i \geq 0 \mid \lg^{(i)} x \leq 1\}. \quad (1.4.30)$$

- Example

$$\begin{aligned} \lg^* 2 &= 1, \\ \lg^* 4 &= 2, \\ \lg^* 16 &= 3, \\ \lg^* 65536 &= 4, \\ \lg^* 2^{65536} &= 5. \end{aligned}$$

The iterative logarithm function is a very slow growing function.

Mathematical Backgrounds, X

Fibonacci numbers

- The **Fibonacci numbers** are defined as

$$\begin{aligned}F_0 &= 0, \\F_1 &= 1, \\F_i &= F_{i-1} + F_{i-2} \quad \text{for } i \geq 2.\end{aligned}\tag{1.4.31}$$

- The first few Fibonacci numbers are

$$0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, \dots$$

- The Fibonacci number is related to the **golden ratio**, ϕ , and its conjugate, $\hat{\phi}$, as

$$F_i = \frac{\phi^i - \hat{\phi}^i}{\sqrt{5}}.\tag{1.4.32}$$

And,

$$\begin{aligned}\phi &= \frac{1 + \sqrt{5}}{2} = 1.61803\dots, \\ \hat{\phi} &= \frac{1 - \sqrt{5}}{2} = -0.61803\dots.\end{aligned}\tag{1.4.33}$$

- It can be shown that

$$F_i = \left\lfloor \frac{\phi^i}{\sqrt{5}} \right\rfloor.\tag{1.4.34}$$

Thus, the Fibonacci numbers grow exponentially.