

# EE3980 Algorithms

## HW6 Linear Sort

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2018/4/15

### Introduction

In this assignment, we're asked to sort a list of words using algorithm of linear time complexity. However, compared with previous sorting assignments. The words to be sorted share two properties, i.e.

1. All words consist of lower-case letters only.
2. The maximum number of letters of the words is 14.

### Approach

Since the characters are all lower-case, which means there's only 27(a ~z and '\0') possible value for each letter in a word string. Plus, the words are no longer than 14 characters (limited length). In such case, a linear-complexity algorithm, radix sort, can be applied.

### Radix Sort

```
1. Algorithm RadixSort(list, N) {  
2.     For i = LSB to MSB do CountingSort(list, N, I);  
3. }
```

RadixSort is simply calling CountingSort from Least Significant Bit (letter) to Most Significant Bit (letter).

It's noteworthy that as we use `scanf` to import data, the characters fill from index 0 (MSB). Then, if the word is shorter than the length of given array, remaining elements in array would be filled with `'\0'`.

## Counting Sort

```
1. Algorithm CountingSort(list, N) {
2.     Init count = { 0, 0, ...0 };
3.     //count has k members, k is all
       possible value in list
4.     for i = list[1] to list[N] do count[i]++;
5.     for i = 2 to k do count[i] += count[i - 1];
6.     for i = N to 1 do A[ --count[ list[i] ] ] = list[i];
7.     return A;
8. }
```

In the above algorithm, first we use `count` array to calculate how many members are less than or equal to the `i`-th possible value. Then, from back to top we place the elements in `list` to `A` according to the position indicated by `count` array.

As we can observe from the looping bounds, the time complexity is

$O(n + k)$ . Where  $n$  is the task size and  $k$  is the number of possible value in `list`.

Additionally, we used another `A` and `count` array, so the space complexity is also

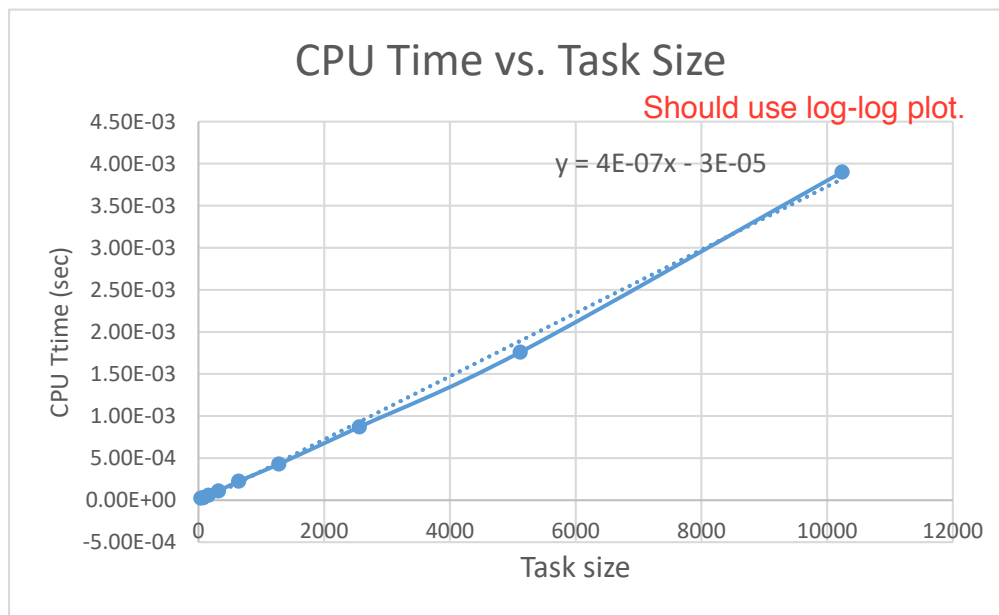
$O(n + k)$ . Therefore the complexity of RadixSort is  $O( r(n + k) )$ .  $r$  is the

maximum length of word in wordlist to be sorted.

## Results and analysis

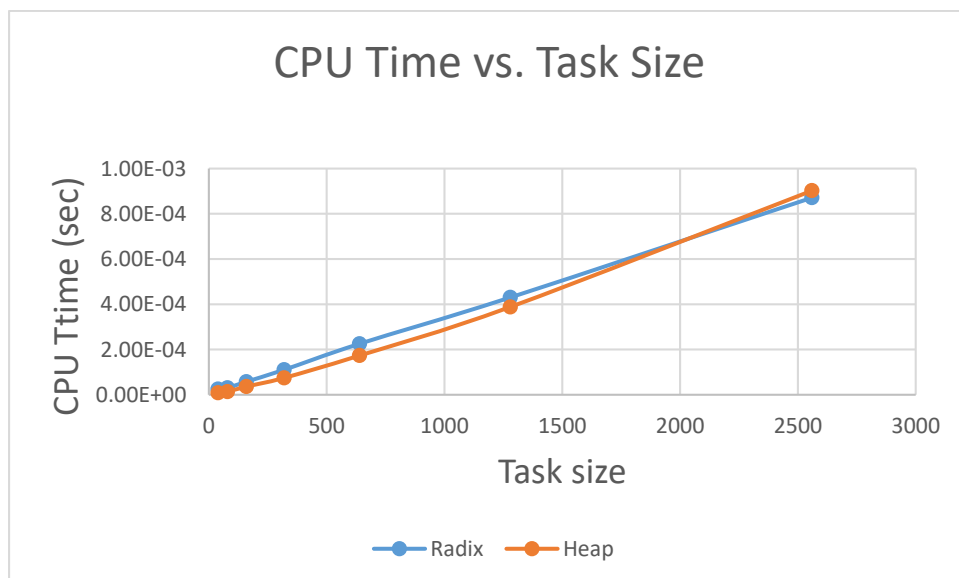
Table. CPU Time (in sec) w.r.t. task size

Task Size	40	80	160	320	640	1280	2560	5120	10240
CPU Time	2.43E-05	3.03E-05	5.75E-05	1.10E-04	2.25E-04	4.30E-04	8.71E-04	1.76E-03	3.90E-03



It's obvious in above chart that when  $r, k \ll n$ , RadixSort has linear time complexity.

However, we can take HeapSort from HW2 to compare together.



Though the theoretical time complexity is different, their actual execution time didn't differ a lot when sorting the test cases of this assignment.

## **Observations and Conclusion**

1. RadixSort / CountingSort are of great use when the data to be sorted have limited possible value. ( $r, k \ll n$ )
2. Lower time complexity does not always guarantee shorter execution time.

```

1 /*****
2   EE3980 Algorithms HW05 Linear Sort
3   Li-Yu Feng 104061212
4   Date:2018/4/15
5 *****/
6
7
8 #include <stdio.h>
9 #include <stdlib.h>
10 #include <string.h>
11 #include <sys/time.h>
12 #include<stdbool.h>
13 #define LEN 14
14 char **A;
15
16
17 double GetTime(void);
18 void RadixSort(char **list,int n);
19 void CountingSort(char **list,int n, int loc);
20 void Heapify(char **list, int i, int n);
21 void HeapSort(char **list,int n);
22 double GetTime(void)
23 {
24     struct timeval tv;
25     gettimeofday(&tv,NULL);
26     return tv.tv_sec+1e-6*tv.tv_usec;
27 }
28
29
30 void CountingSort(char **list, int n, int loc){
31     int count[27];           //n:size loc:sorting radix location
32     int i,index;
33     char *temp;
34
35
36     for (i = 0; i < 27; ++i)    //init count array
37     {
38         count[i] = 0;
39     }
40
41
42     for (i = 0; i < n; i++){    //count alphabet occuring times
43         if(list[i][loc] == '\0')
44             count[0]++;
45         else count[ list[i][loc] - 96 ]++;
46     }
47
48
49     for(i = 1; i < 27; i++){    // calculate each alphabet's
50         count[i] += count[i-1]; // dictionary order

```

Need comments here to explain the purpose of each function.

```

51     }
52
53
54     for(i = n-1; i >= 0; i-- ){
55         temp = list[i];
56         if (temp[loc] == '\0')    index = 0;    //deal with '\0' (ASCII 0)
57         else index = temp[loc] - 96;    //lower case alphabet
58         A[ count[ index ]- 1] = temp;    //put word in new array
59         count[ index ]--;    //update count array
60     }
61
62     for ( i = 0; i < n; ++i)    //output current result
63     {
64         list[i] = A[i];
65     }
66 }
67 }
68
69
70
71 void RadixSort(char **list,int n){
72     int i;
73
74     for (i = LEN-1; i >= 0; --i){
75         CountingSort(list, n, i);
76     }
77
78 }
79
80 void Heapify(char **list, int i, int n){
81     int j = i*2;
82     char *temp = list[i-1];
83     bool done = false;
84
85     while(j<=n && !done){
86         if(j<n && strcmp(list[j-1],list[j+1-1]) < 0) j++;
87         if(strcmp(temp,list[j-1] ) > 0 ) done = true;
88         else{
89             list[j/2-1] = list[j-1];
90             j *= 2;
91         }
92         //printf("%d\n",j);
93     }
94     list[j/2-1] = temp;
95 }
96
97 void HeapSort(char **list,int n){
98     char *temp;
99     int i;
100

```

```

101     for( i = n/2 ; i>0 ; i--)
102         {Heapify(list,i,n);}
103     for(i = n; i > 1; i-- ){
104         temp = list[i-1];
105         list[i-1] = list[0];
106         list[0] = temp;
107         Heapify(list,1,i-1);
108     }
109 }
110
111 int main(){
112     int i,j;
113     int Nwords;
114     double t;
115     char **words,**temp;
116
117     scanf("%d", &Nwords);
118     A = (char **)malloc( Nwords * sizeof(char*) );
119     temp = (char **)malloc( Nwords * sizeof(char*) );
120     words = (char**)malloc(Nwords * sizeof(char*)); //
121     for(i = 0; i < Nwords; i++) //
122         words[i] = (char *)malloc( (LEN+1) * sizeof(char)); //
123
124
125     for(i = 0; i < Nwords ; i++){ //
126         scanf("%s", words[i]); //scan words
127     }
128     t = GetTime();
129
130     for(i =0; i<500;i++){ //sort 500 times
131         for (j = 0; j < Nwords; ++j)
132             {
133                 temp[j] = words[j];
134             }
135         RadixSort(temp, Nwords);
136     }
137
138
139     for(i = 0; i < Nwords ; i++){ //print result
140         printf("%d %s\n",i, temp[i]);
141     }
142
143     t = GetTime() - t;
144     printf("%s:\nN=%d\nCPU time = %.3g seconds\n", "Linear Sort",
145           Nwords, t / 500.0);
146
147     t = GetTime();
148
149     for(i =0; i<500;i++){ //sort 500 times
150         for (j = 0; j < Nwords; ++j)

```

```

151     {
152         temp[j] = words[j];
153     }
154     HeapSort(temp, Nwords);
155 }
156
157
158     for(i = 0; i < Nwords ; i++){           //print result
159         printf("%d %s\n",i, temp[i]);
160     }
161
162     t = GetTime() - t;
163     printf("%s:\nN=%d\nCPU time = %.3g seconds\n", "Heap Sort",
164           Nwords, t / 500.0);
165
166
167     return 0;
168
169 }

```



Score: 91

---

[Table] can be better presented.

[Figure] can be better presented.

- Use log-log scale.
- Can compare two algorithms on the same plot.